

# Some Elliptic CURve REsults

## Elligator

by Daniel J. Bernstein, Mike Hamburg, Anna Krasnova,  
and Tanja Lange

## MiniaLT

by W. Michael Petullo, Xu Zhang, Jon A. Solworth,  
Daniel J. Bernstein, and Tanja Lange

<http://safecurves.cr.yp.to/>

by Daniel J. Bernstein and Tanja Lange

December 3, 2013

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Your protocol sends elliptic curve points (public keys) but they need to look like random strings because

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You might be interested in Elligator!

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Choose prime  $p = 2^n - c$  with small  $c$ , then integers  $\leq (p - 1)/2$  cover essentially all strings of  $n - 1$  bits.

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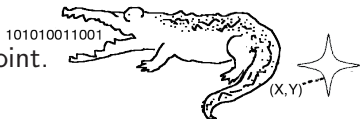
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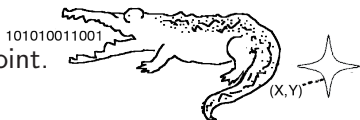
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Maps are inverse of one another and efficiently computable.  
Very fast test whether inverse map is defined on point  
 $\Rightarrow$  fast rejection of points, tweak protocol to try again.

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# MinimaLT – Minimal Latency Tunneling

- ▶ MinimaLT protocol makes secure connections
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# MinimalT – Minimal Latency Tunneling

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- ▶ Skips TCP handshake; authentication makes this safe
- ▶ Provides IP-address mobility: unlinkable, no handshake
- ▶ Overlaps ephemeral-key retrieval with DNS lookup
- ▶ Erases keys within minutes (TLS often takes months)
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Extra advertisement

small (human-auditable) codebase: [TweetNaCl](#)

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<http://www.ethos-os.org/>

W. Michael Petullo, Xu Zhang, Jon A. Solworth, Daniel J. Bernstein, & me

# SafeCurves: choosing safe curves for elliptic-curve cryptography

All known security criteria for elliptic curves,  
machine verified.

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Also: can the curve be backdoored?

# SafeCurves: choosing safe curves for elliptic-curve cryptography

Curve	Safe?	Parameters:			ECDLP security:				ECC security:			
		field	equation	base	rho	transfer	disc	rigid	ladder	twist	complete	ind
Anomalous	False	True ✓	True ✓	True ✓	True ✓	False	False	True ✓	False	False	False	False
M-221	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓
E-222	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓
NIST P-224	False	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	False	False	False	False
Curve1174	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓
Curve25519	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓
BN(2,254)	False	True ✓	True ✓	True ✓	True ✓	False	False	True ✓	False	False	False	False
brainpoolP256t1	False	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	False	False	False	False
ANSSI FRP256v1	False	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	False	False	False	False
NIST P-256	False	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	False	False	True ✓	False
secp256k1	False	True ✓	True ✓	True ✓	True ✓	True ✓	False	True ✓	False	True ✓	False	False
M-383	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓
Curve383187	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓
brainpoolP384t1	False	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	False	True ✓	False	False
NIST P-384	False	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	False	False	True ✓	False
Curve3617	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓	True ✓

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